



ComPAS
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Neuchâtel

**Understanding and managing
hardware affinities
on hierarchical platforms
With Hardware Locality (hwloc)**

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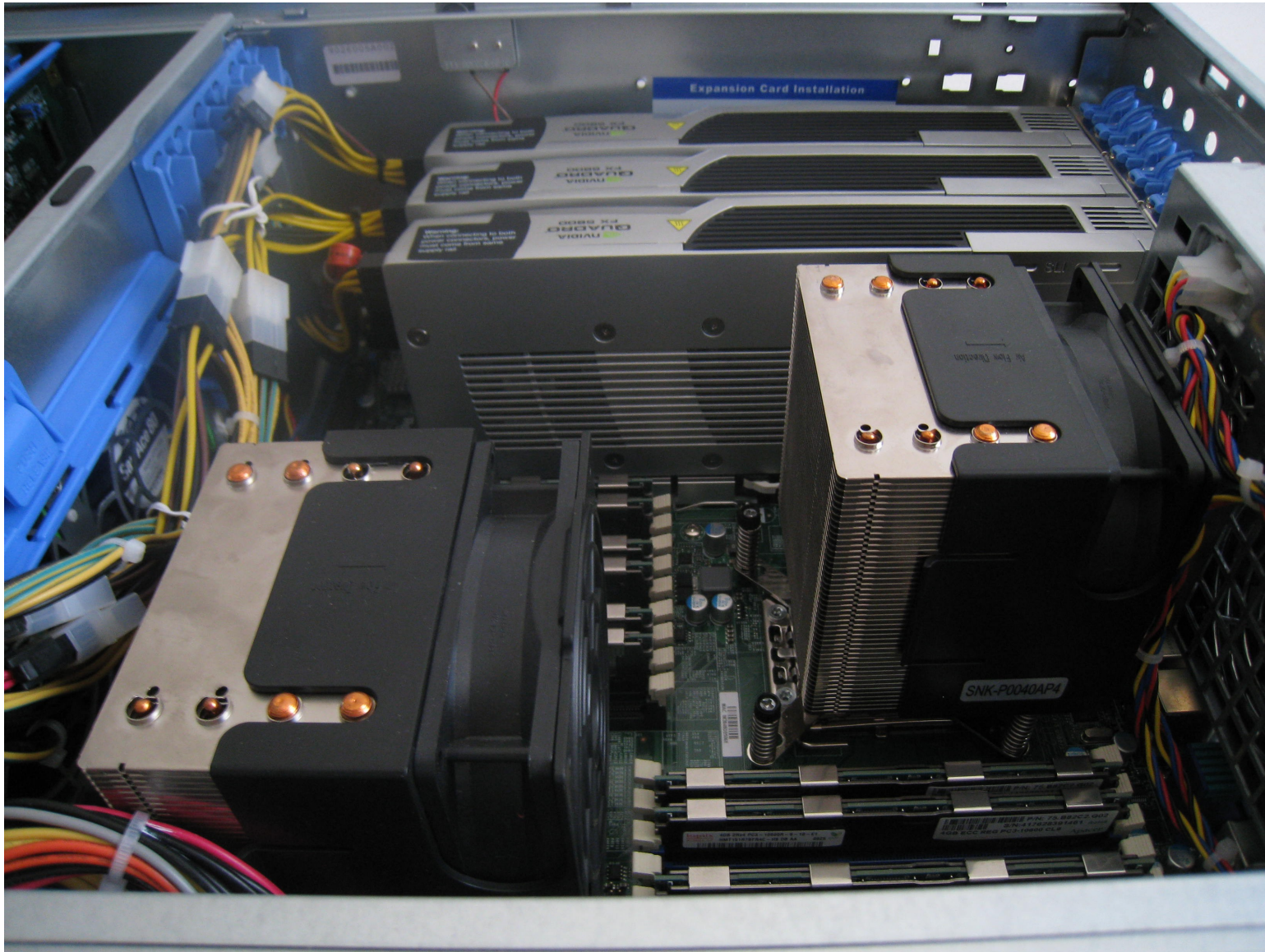
Agenda

- Quick example as an Introduction
- Bind your processes
- What's the actual problem?
- Introducing hwloc (Hardware Locality)
- Command-line tools
- C Programming API
- Use cases
- Conclusion

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**Quick example
as an Introduction**

Machines are increasingly complex



Machines are increasingly complex

- Multiple processor sockets
- Multicore processors
- Simultaneous multithreading
- Shared caches
- NUMA
- GPUs, NICs, ...
 - Close to some sockets (NUIOA)

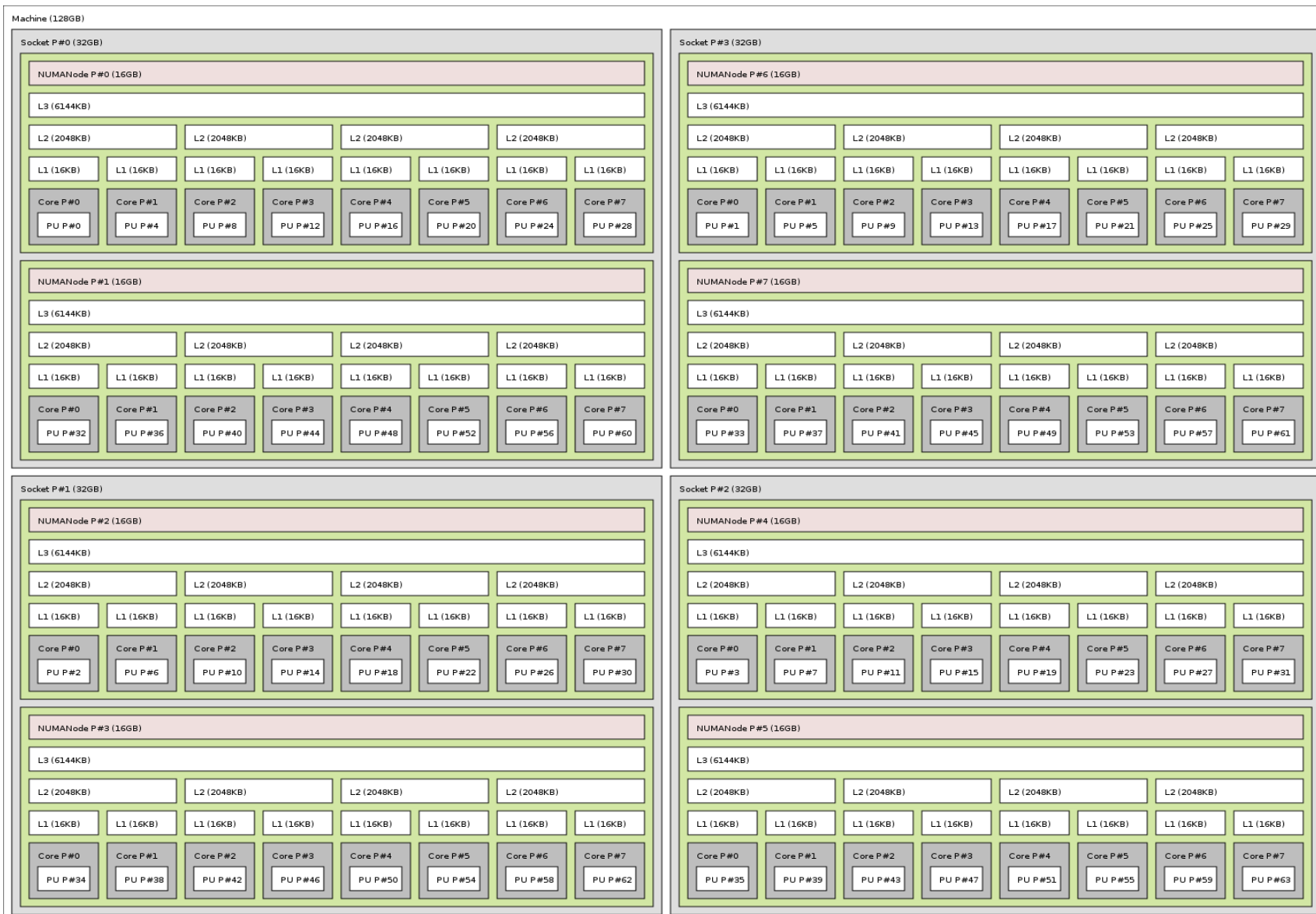
Example with MPI

- Let's say I have a 64-core AMD machine
 - Not unusual (about 6000\$)
- I am running a MPI pingpong between pairs of cores
 - Open MPI 1.6
 - Intel MPI Benchmarks 3.2

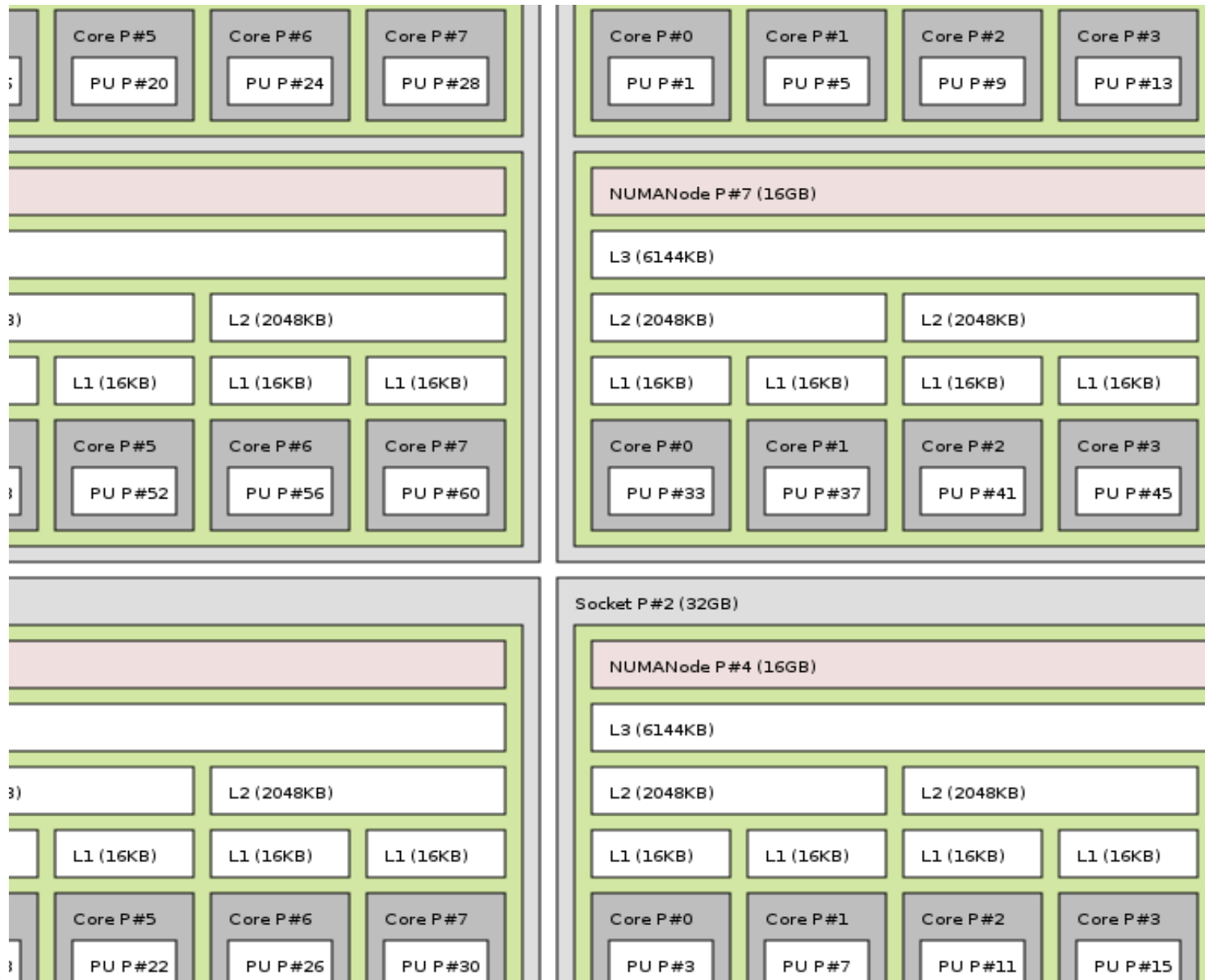
Example with MPI (2/3)

- Between cores 0 and 1
 - 1.39 μs latency – 1900MB/s throughput
- Between cores 0 and 4
 - 1.63 μs – 1400 MB/s – Interesting !
- Between cores 0 and 5
 - 0.68 μs – 3600 MB/s – What ?!
- Between cores 0 and 8
 - 1.24 μs – 2400 MB/s
- Between cores 0 and 32
 - 1.34 μs – 2100 MB/s

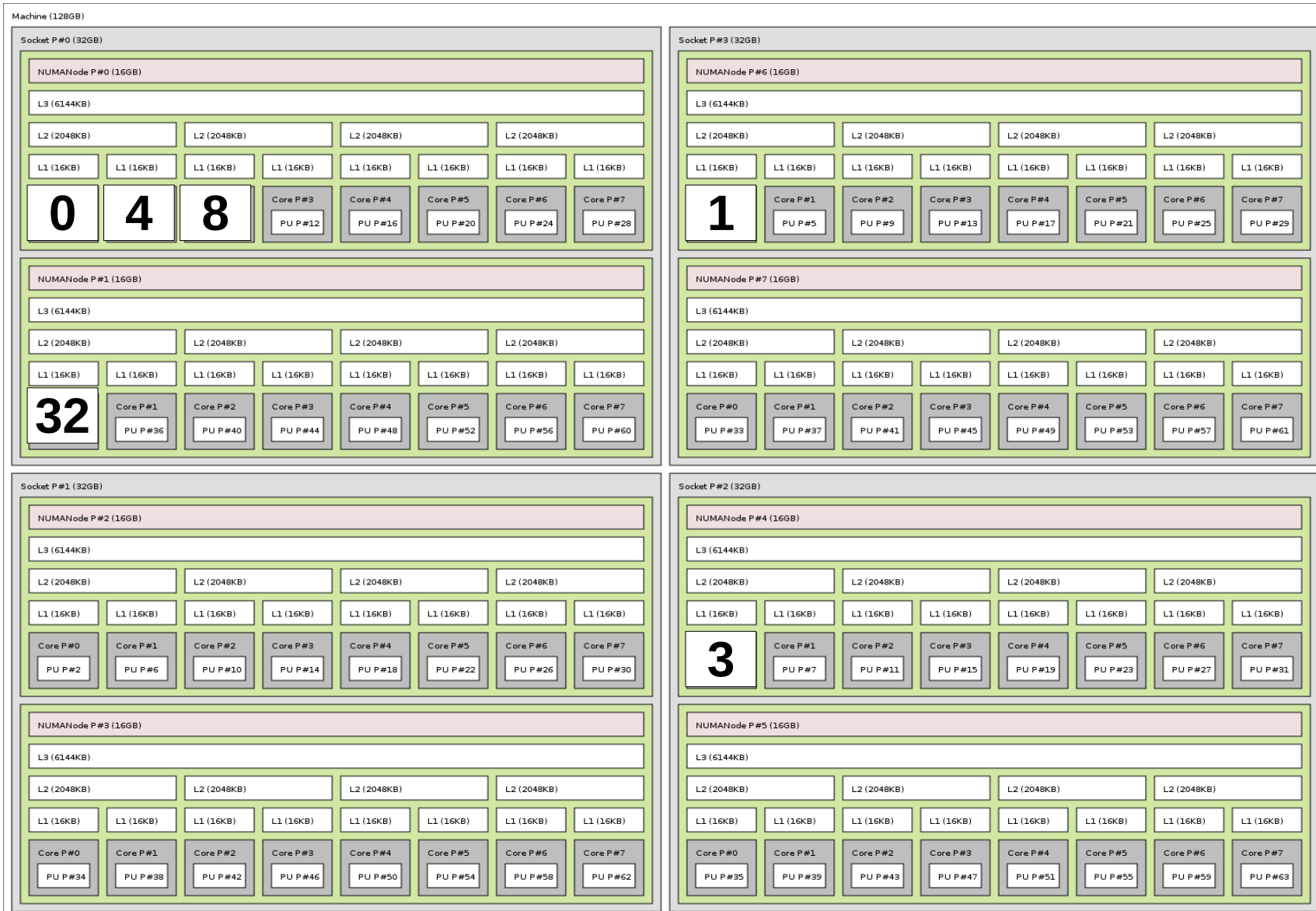
What is going on



What is going on (2/3)



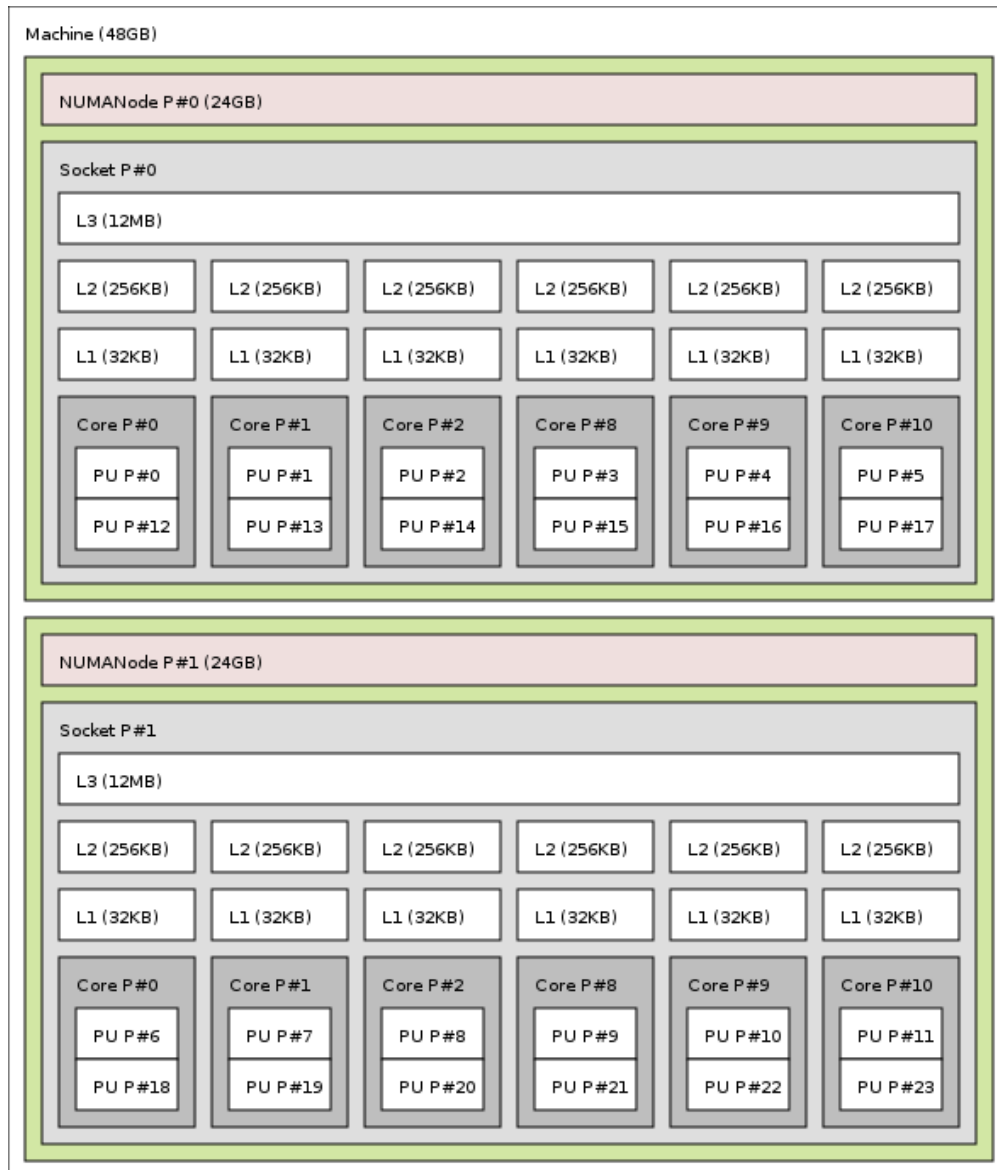
What is going on (3/3)



Example with MPI (3/3)

- Between cores that share a L2 cache
 - 0.68 μs – 3600 MB/s
- Between cores that only share a L3 cache
 - 1.24 μs – 2400 MB/s
- Between cores inside the same socket
 - 1.34 μs – 2100 MB/s
- Between cores of another socket
 - 1.39 μs – 1900MB/s
- Between cores of another socket further away
 - 1.63 μs – 1400 MB/s

Ok, what about Intel machines?



- Less hierarchy levels
 - 4 vs 3
 - HyperThreading?
- But same problems

First take away messages

- Locality matters to communication performance
 - Machines are really far from flat
- Cores/processors numbering is crazy
 - Never expect anything sane here

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Bind your processes

Where does locality actually matter?

- MPI communication between processes on the same node
- Shared-memory too (threads, OpenMP, etc)
 - Synchronization
 - Barriers use caches and memory too
 - Concurrent access to shared buffers
 - Producer-consumer, etc
- 10 years ago, locality was mostly an issue for large NUMA SMP machines (SGI, etc)
 - Today it's **everywhere**
 - Because multicores and NUMA are everywhere

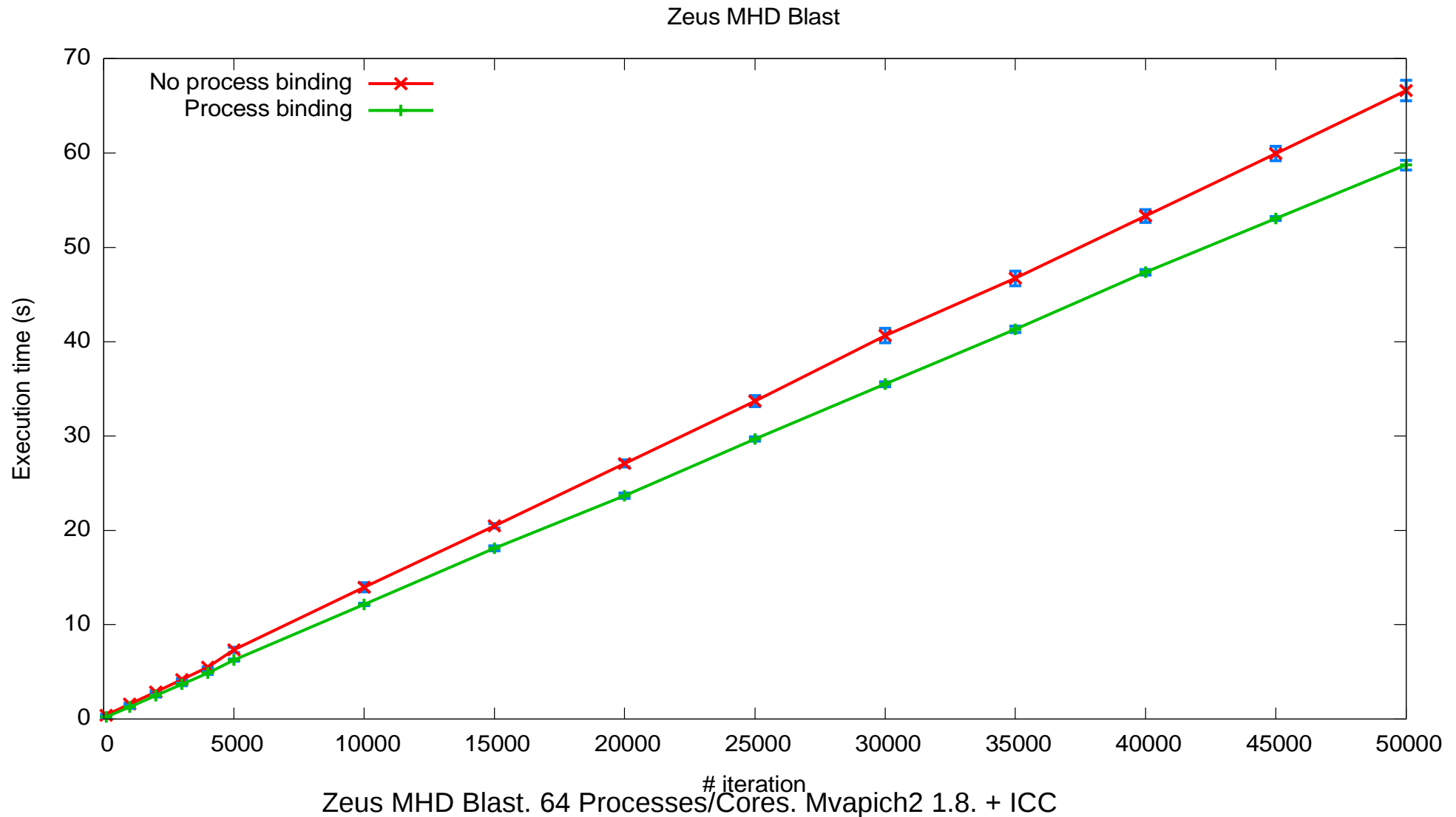
What to do about locality?

- Place processes/tasks according to their affinities
 - If two tasks communicate/synchronize/share a lot, keep them close
- Adapt your algorithms to the locality
 - Adapt communication/synchronization implementations to the topology
 - Ex: hierarchical barriers

Process binding

- Some MPI implementations bind processes by default (Intel MPI, Open MPI 1.8)
 - Because it's better for reproducibility
- Some don't
 - Because it may hurt your application
 - Oversubscribing?
- Binding doesn't guarantee that your processes are optimally placed
 - It just means your process won't move
 - No migration, less cache issues, etc

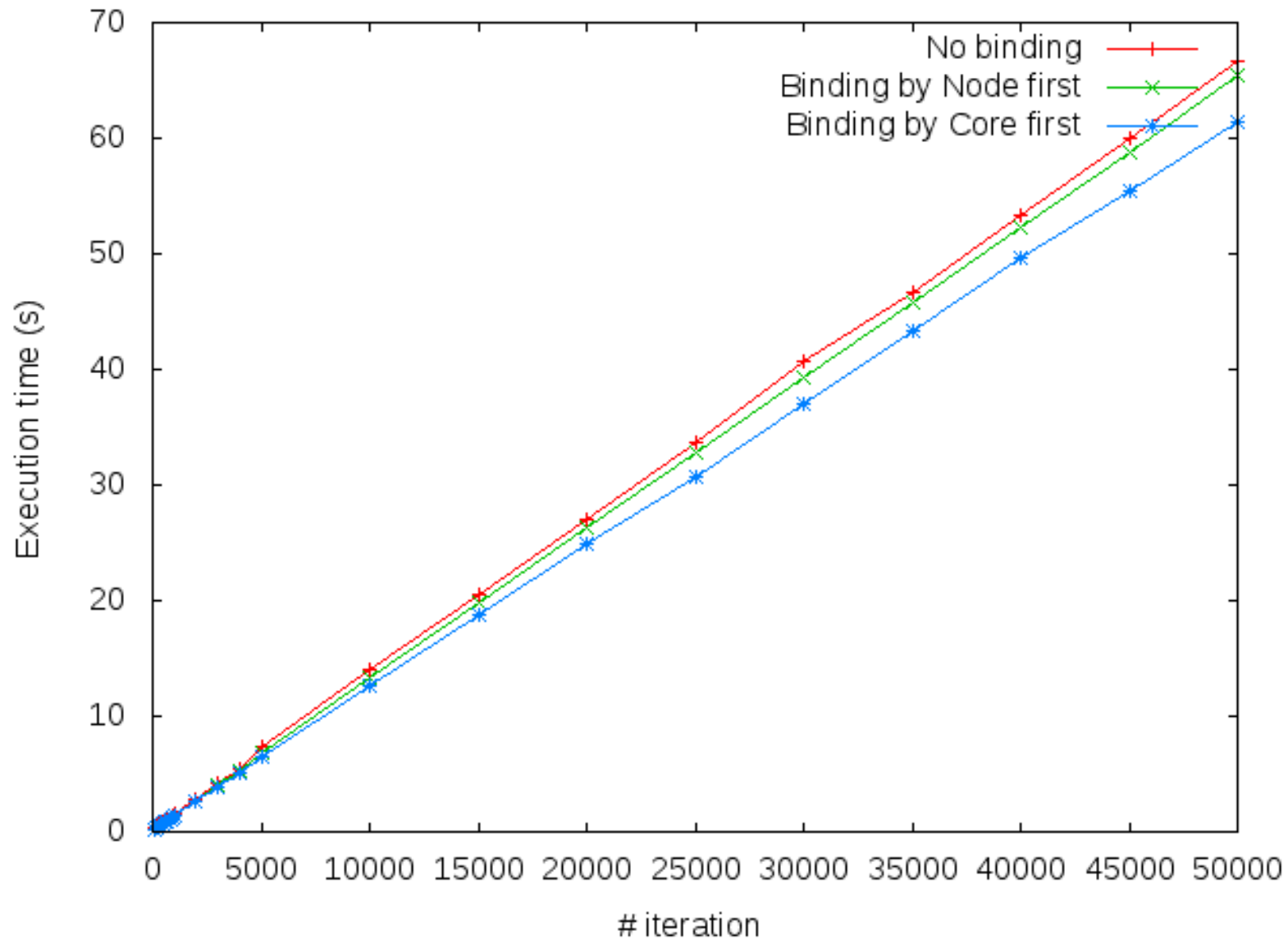
To bind or not to bind ?



Where to bind ?

- Default binding strategies ?
 - By core first :
 - One process per core on first node, then one process per on second node, ...
 - By node first :
 - One process on first core of each node, then one process on second core on each node, ...
- Your application likely prefers one to the other
 - Usually the first one
 - Because you often communicate with nearby ranks

Binding strategy impact



How to bind in MPI? (more later)

- MPI standard says nothing
- Open MPI
 - `mpiexec --bind-to core -np 8 -H node1,node2 ./myprogram`
- MPICH
 - `mpiexec -bind-to core ...`
- Manually
 - `mpiexec`
 - np 1 -H node1 numactl --physcpubind 0 ./myprogram :
 - np 1 -H node1 numactl --physcpubind 1 ./myprogram :
 - np 1 -H node2 numactl --physcpubind 0 ./myprogram
 - Rank files, etc

How to bind in OpenMP? (more later)

- Intel Compiler

- KMP_AFFINITY=scatter or compact

- GCC

- GOMP_CPU_AFFINITY=1,3,5,2,4,6

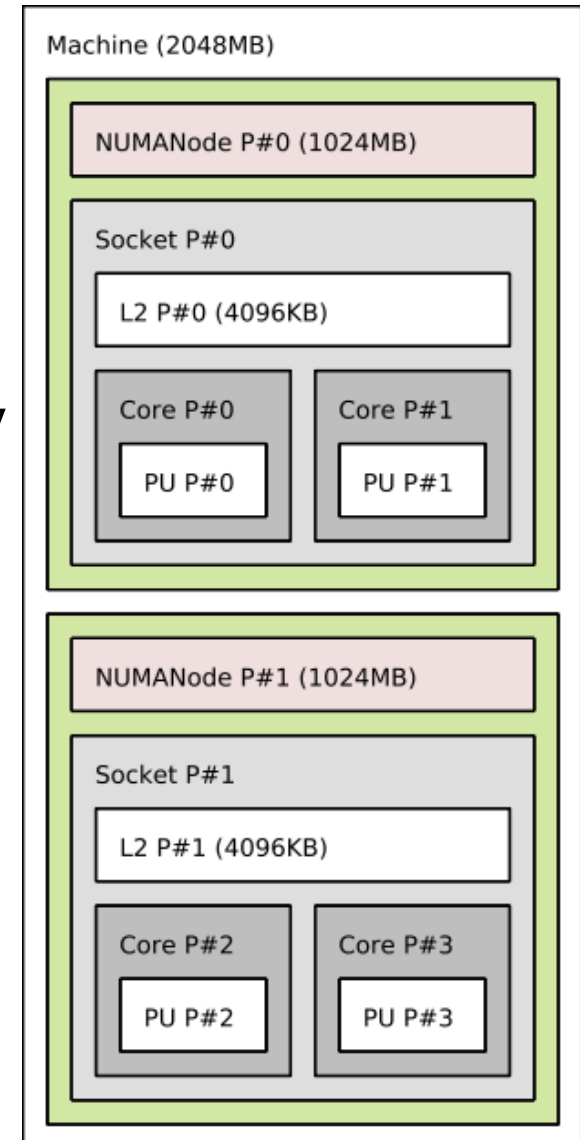
How do I choose?

- Dilemma

- Use cores 0 & 1 to share cache and improve synchronization cost ?
- Use core 0 & 2 to maximize memory bandwidth ?

- Depends on

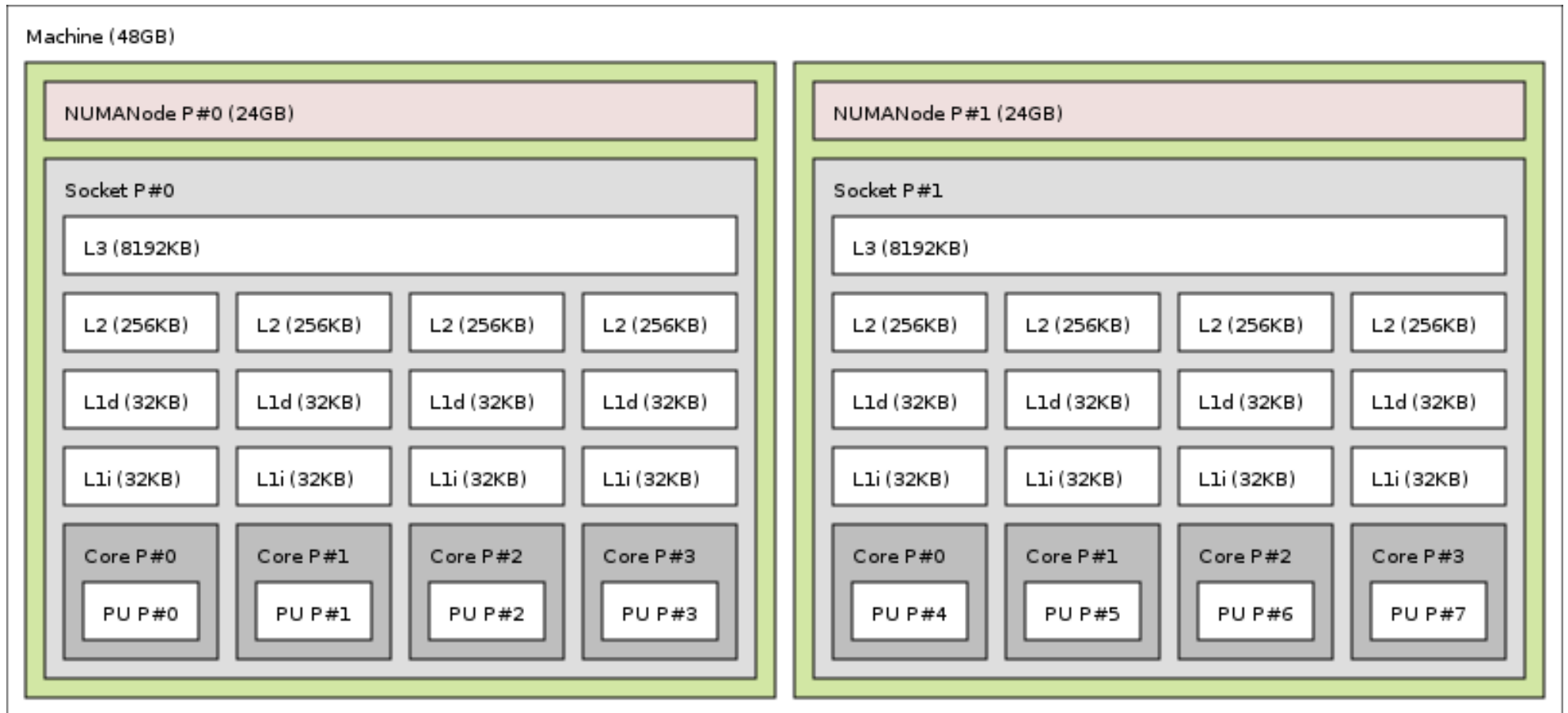
- The machine structure
- The application needs



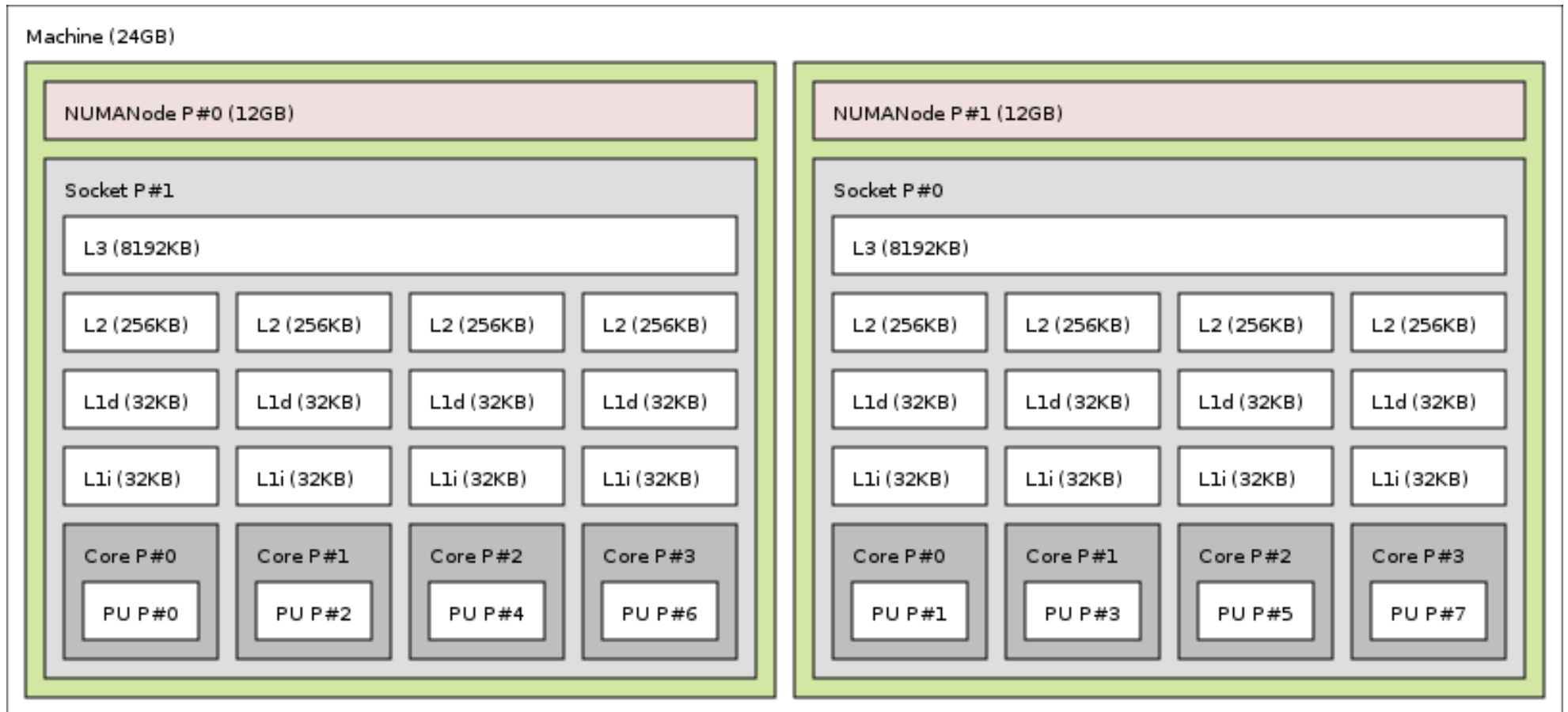
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What's the actual
problem ?

Example of dual Nehalem Xeon machine



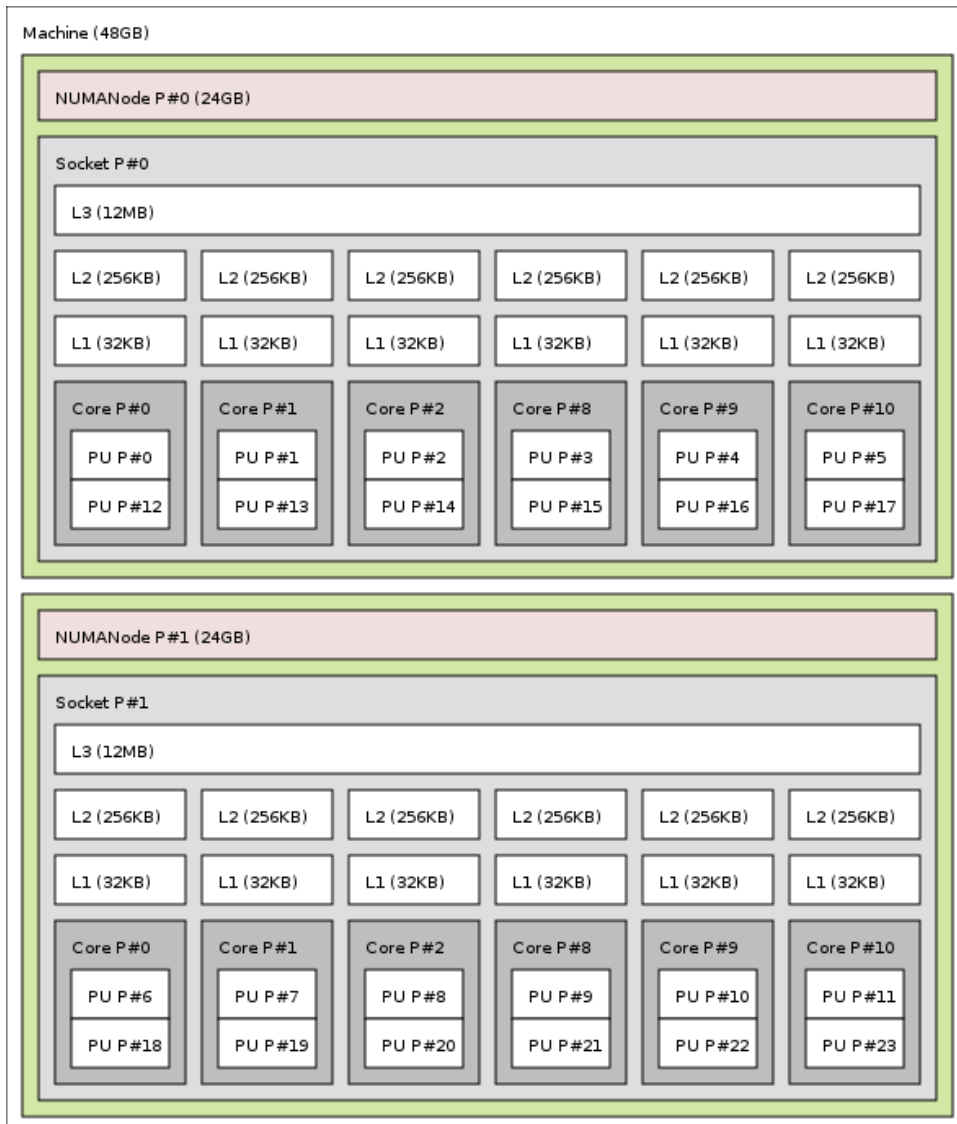
Another example of dual Nehalem Xeon machine



Processor and core numbers are crazy

- Resources ordering is unpredictable
 - Ordered by any combination of NUMA/socket/core/hyperthread
 - Can change with the vendor, the BIOS version, etc
- Some resources may be unavailable
 - Batch schedulers can give only parts of machines
 - Core numbers may be non-consecutive, non starting at 0, etc
- Don't assume anything about indexes
 - Don't use these indexes
 - Or you won't be portable

Level ordering isn't much better



- Intel is usually
 - Machine
 - Socket = NUMA = L3
 - Core = L1 = L2
 - Hyperthread (PU)

Level ordering isn't much better (2/3)



- AMD is different

- Machine
- Socket
- NUMA = L3
- L2 = L1i
- Core = L1d

Level ordering isn't much better (3/3)

- Sometimes there are multiple sockets per NUMA nodes
 - And different levels of caches
- Don't assume anything about level ordering
 - Or (again) you won't be portable
 - e.g.: Intel Compiler OpenMP binding may be wrong on AMD machines

Gathering topology information is difficult

- Lack of generic, uniform interface
 - Operating system specific
 - /proc and /sys on Linux
 - rset, sysctl, lgrp, kstat on others
 - Hardware specific
 - x86 cpuid instruction, device-tree, PCI config space, ...
- Evolving technology
 - AMD Bulldozer dual-core compute units
 - It's not two real cores, neither a dual-threaded core
 - New levels? New ordering?

Binding is difficult too

- Lack of generic, uniform interface, again
 - Process/thread binding
 - sched_setaffinity API changed twice on Linux
 - rset, ldom_bind, radset, affinity_set on others
 - Memory binding
 - mbind, migrate_pages, move_pages on Linux
 - rset, mmap, radset, nmadvise, affinity_set on others
 - Different constraints
 - Bind on single core only, on contiguous set of cores, on random sets ?
 - Many different policies

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Introducing hwloc (Hardware Locality)

What hwloc is

- Detection of hardware resources
 - Processing units (PU), logical processors, hardware threads
 - Everything that can run a task
 - Memory nodes, shared caches
 - Cores, Sockets, ... (things that contain multiple PUs)
 - I/O devices
 - PCI devices and corresponding software handles
- Described as a tree
 - Logical resource identification and organization
 - Based on locality

What hwloc is (2/2)

- API and tools to consult the topology
 - Which cores are near this memory node ?
 - Give me a single thread in this socket
 - Which memory node is near this GPU ?
 - What shared cache size between these cores ?
- Without caring about hardware strangeness
 - Non portable and crazy numbers, names, ...
- A portable binding API
 - No more Linux sched_setaffinity API breakage
 - No more tens of different binding API with different types

What hwloc is not

- A placement algorithm
 - hwloc gives hardware information
 - You're the one that knows what your software does/needs
 - You're the one that must match software affinities to hardware localities
 - We give you the hardware information you need
- A profiling tool
 - Other tools (e.g. likwid) give you hardware performance counters
 - hwloc can match them with the actual resource organization

History

- Runtime Inria project in Bordeaux, France
 - Thread scheduling over NUMA machines (2003...)
 - Marcel threads, ForestGOMP OpenMP runtime
 - Portable detection of NUMA nodes, cores and threads
 - Linux wasn't that popular on NUMA platforms 10 years ago
 - Other Unixes have good NUMA support
 - Extended to caches, sockets, ... (2007)
 - Raised questions for new topology users
 - MPI process placement (2008)

History

- Marcel's topology detection extracted as standalone library (2009)
- Noticed by the Open MPI community
 - They knew their PLPA library wasn't that good
- Merged both libraries as hwloc (2009)
- BSD-3
- Still mainly developed by Inria Bordeaux
 - Collaboration with Open MPI community
 - Contributions from MPICH, Redhat, IBM, Oracle, ...

Alternative software with advanced topology knowledge

- PLPA (old Open MPI library)
 - Linux specific, no NUMA support, obsolete, dead
- libtopology (IBM)
 - Dead
- Likwid
 - x86 only, needs update for each new processor generation, no extensive C API
 - It's more kind of a performance optimization tool
- Intel Compiler (icc)
 - x86 specific, no API

hwloc's view of the hardware

- Tree of objects
 - Machines, NUMA memory nodes, sockets, caches, cores, threads
 - Logically ordered
 - Grouping similar objects using distances between them
 - Avoids enormous flat topologies
 - Many attributes
 - Memory node size
 - Cache type, size, line size, associativity
 - Physical ordering
 - Miscellaneous info, customizable

Installing hwloc

- Packages available in Debian, Ubuntu, Redhat, Fedora, CentOS, ArchLinux, NetBSD
- You want the development headers too
 - libhwloc-dev, hwloc-devel, ...

Manual installation

- Take a recent tarball at <http://www.open-mpi.org/projects/hwloc>
- Dependencies
 - On Linux, numactl/libnuma development headers
 - Cairo headers for Istopo graphics
- `./configure --prefix=$PWD/install`
 - Very few configure options
- Check the summary at the end of configure

Manual installation

- make
- make install
- Useful environment variables
 - export PATH=\$PATH:<prefix>/bin
 - export LD_LIBRARY_PATH=\$LD_LIBRARY_PATH:<prefix>/lib
 - export
PKG_CONFIG_PATH=\$PKG_CONFIG_PATH:<prefix>/lib/pkgconfig
 - export MANPATH=\$MANPATH:<prefix>/share/man
- Have access to a nice server for this tutorial ?
 - Install hwloc on the server AND on your local machine

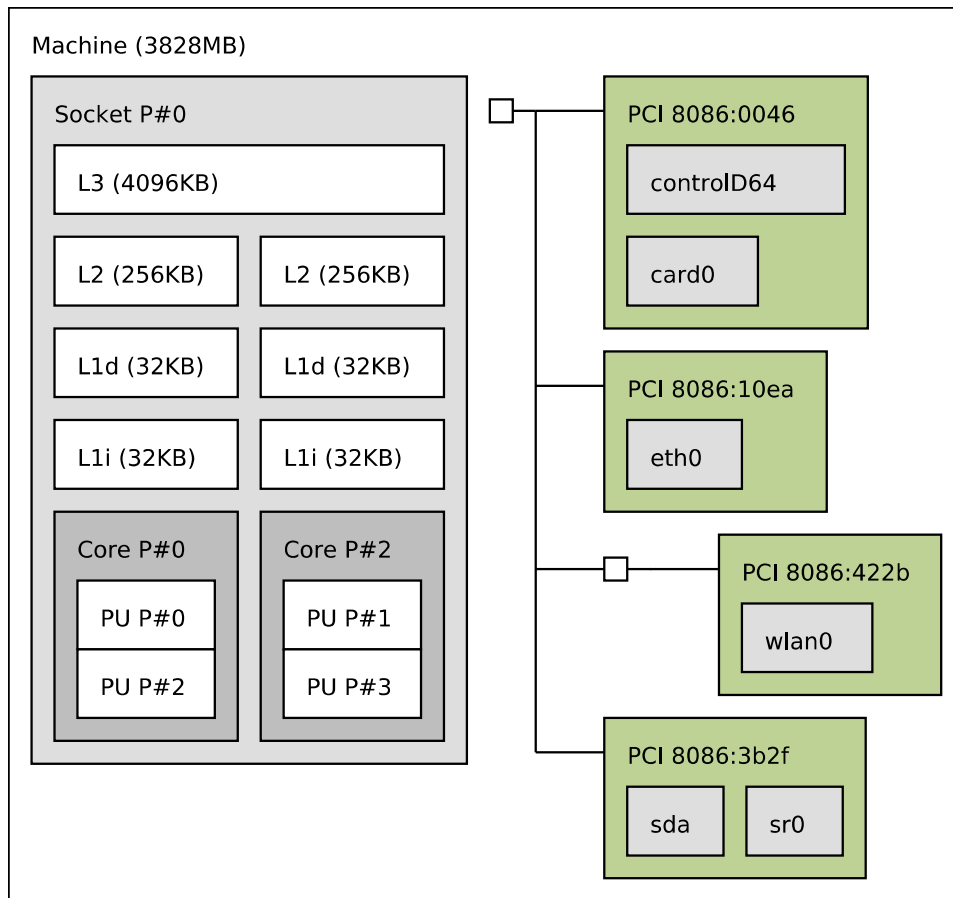
Using hwloc

- Many hwloc command-line tools
 - lstopo and hwloc-*
- ... but the actual hwloc power is in the C API
- Perl and Python bindings

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Command-line Tools

Istopo (displaying topologies)



Machine (3828MB)

Socket L#0 + L3 L#0 (4096KB)

L2 L#0 (256KB) + Core L#0

PU L#0 (P#0)

PU L#1 (P#2)

L2 L#1 (256KB) + Core L#1

PU L#2 (P#1)

PU L#3 (P#3)

HostBridge L#0

PCI 8086:0046

GPU L#0 "controlD64"

PCI 8086:10ea

Net L#2 "eth0"

PCIBridge

PCI 8086:422b

Net L#3 "wlan0"

PCI 8086:3b2f

Block L#4 "sda"

Block L#5 "sr0"

Istopo

- Many output formats
 - Text, Cairo (PDF, PNG, SVG, PS), Xfig, ncurses
 - Automatically guessed from the file extension
- XML dump/reload
 - Faster, convenient for remote debugging
- Configuration options for nice figures for papers
 - Horizontal/Vertical placement
 - Legend
 - Ignoring things
 - Creating fake topologies

Istopo

```
$ Istopo
```

```
$ Istopo --no-io -
```

```
$ Istopo myfile.png
```

```
$ ssh host Istopo saved.xml
```

```
$ Istopo -i saved.xml
```

```
$ ssh myhost Istopo -.xml | Istopo --if xml -i -
```

```
$ Istopo -i "node:4 socket:2 core:2 pu:2"
```


hwloc-bind

(binding processes, threads and memory)

- Bind a process to a given set of CPUs

```
$ hwloc-bind socket:1 -- mycommand myargs...
```

```
$ hwloc-bind os=mlx4_0 -- mympiprogram ...
```

- Bind an existing process

```
$ hwloc-bind --pid 1234 node:0
```

- Bind memory

```
$ hwloc-bind --membind node:1 --cpubind node:0 ...
```

- Find out if a process is already bound

```
$ hwloc-bind --get --pid 1234
```

```
$ hwloc-ps
```

hwloc-calc (calculating with objects)

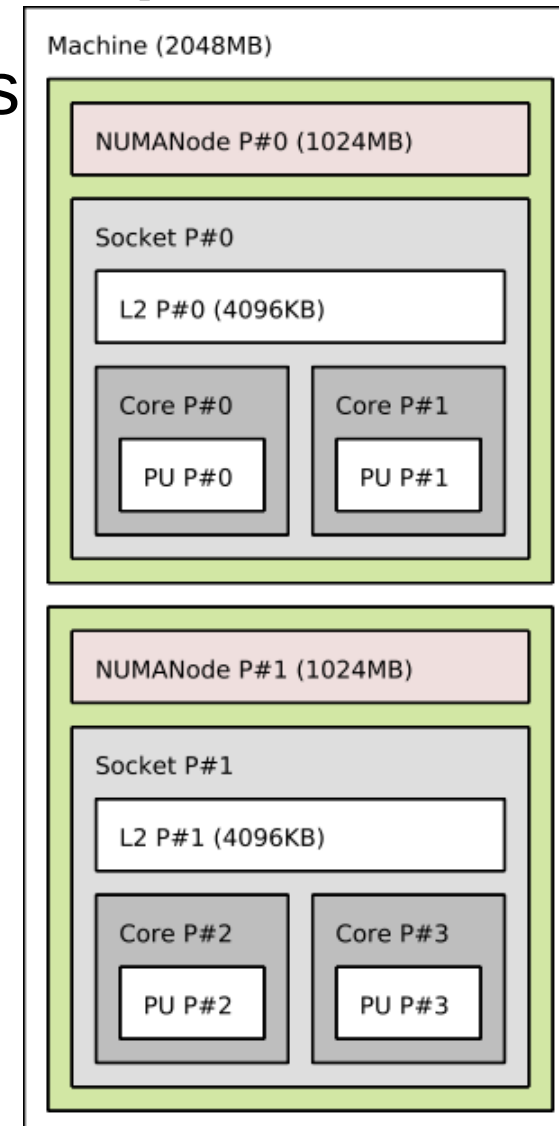
- Convert between ways to designate sets of CPUs, objects... and combine them

```
$ hwloc-calc socket:1.core:1 ~pu:even  
0x00000008
```

```
$ hwloc-calc --number-of core node:0  
2
```

```
$ hwloc-calc --intersect pu socket:1  
2,3
```

- The result may be passed to other tools
 - Multiple invocations may be combined
 - I/O devices also supported
- ```
$ hwloc-calc os=eth0
```



# Other tools

- Get some object information
  - hwloc-info (v1.7+)
- Generate bitmaps for distributing multiple processes on a topology
  - hwloc-distrib
- Save a Linux node topology info for debugging
  - hwloc-gather-topology
- Manipulating multiple topologies, etc.

# 6

## C Programming API

# API basics

- A hwloc program looks like this

```
#include <hwloc.h>
```

```
hwloc_topology_t topo;
```

```
hwloc_topology_init(&topo);
```

```
/* ... configure what topology to build ... */
```

```
hwloc_topology_load(topo);
```

```
/* ... play with the topology ... */
```

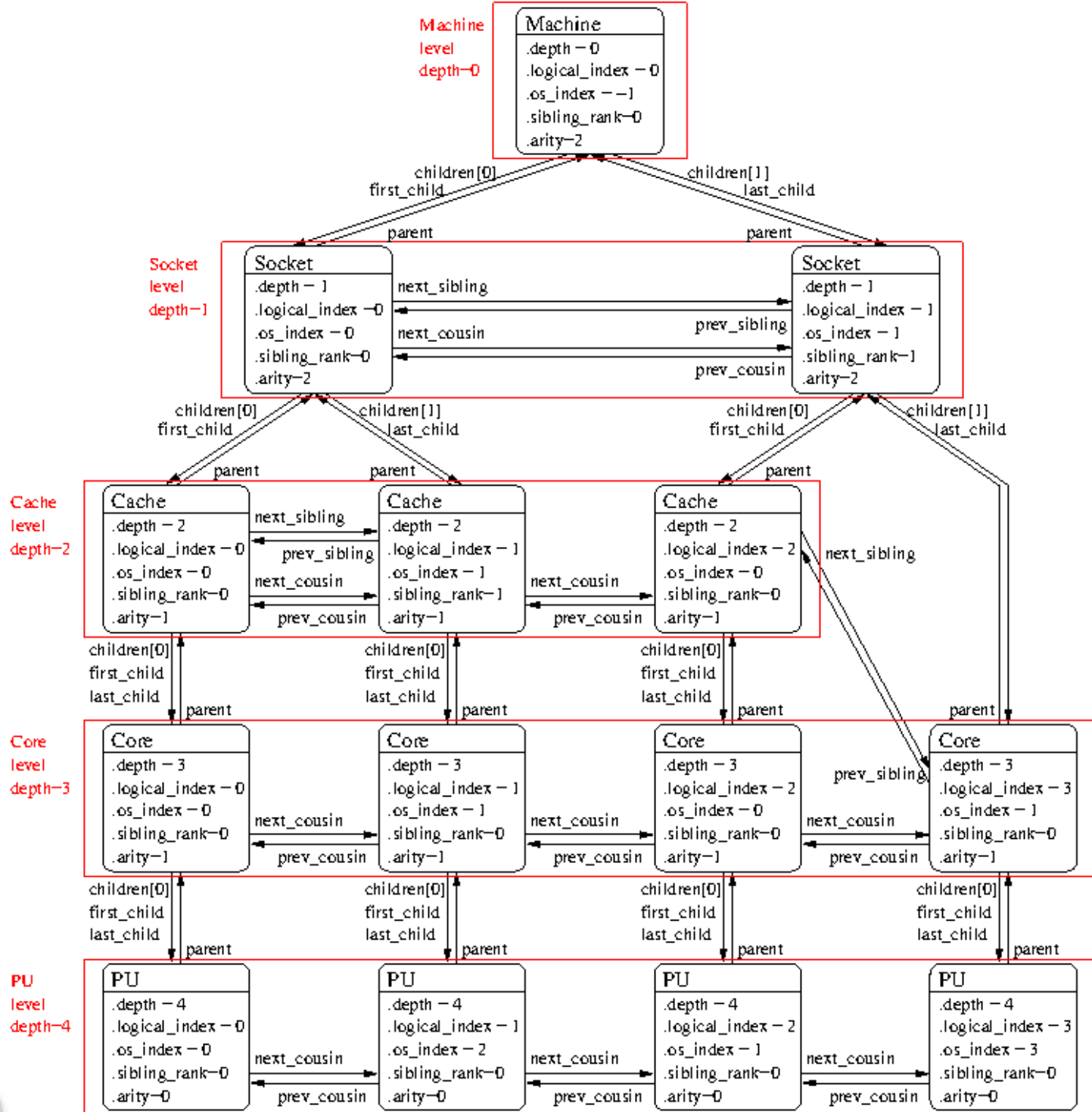
```
hwloc_topology_destroy(topo);
```

# Major hwloc types

- The topology context : `hwloc_topology_t`
  - You always need one
- The main hwloc object : `hwloc_obj_t`
  - That's where the actual info is
  - The structure isn't opaque
    - It contains many pointers to ease traversal
- Object type : `hwloc_obj_type_t`
  - `HWLOC_OBJ_PU`, `_CORE`, `_NODE`, ...

# Object information

- Type
- Optional name string
- Indexes (see later)
- cpusets and nodesets (see later)
- Tree pointers (\*cousin, \*sibling, arity, \*child\*, parent)
- Type-specific attribute union
  - obj->attr->cache.size
  - obj->attr->pcidev.linkspeed
- String info pairs





# Browsing as a tree

- The root is `hwloc_get_root_obj(topo)`
- Objects have children
  - `obj->arity` is the number of children
  - The array of children is `obj->children[]`
  - They are also in a list
    - `obj->first_child`, `obj->last_child`
    - `child->prev_sibling`, `child->next_sibling`
    - NULL-terminated
- The parent is `obj->parent` (or NULL)

# Browsing as levels

- The topology is also organized as levels of identical objects
  - Cores, L2d Caches, ...
  - All PUs at the bottom
- Number of levels `hwloc_topology_get_depth(topo)`
- Number of objects on a level
  - `hwloc_get_nbobjs_by_type(topo, type)`
  - `hwloc_get_nbobjs_by_depth(topo, depth)`
- Convert between depth and type using
  - `hwloc_get_type_depth()` or `hwloc_get_depth_type()`

# Browsing as levels

- Find objects by level and index
  - `hwloc_get_obj_by_type(topo, type, index)`
  - There are variants taking a depth instead of a type
    - Note : the depth of my child is not always my depth + 1
      - Think of asymmetric topologies
- Iterate over objects of a level
  - Objects at the same levels are also interconnect by `prev/next_cousin` pointers
    - Don't mix up siblings (children list) and cousins (level)
  - `hwloc_get_next_obj_by_type/depth()`

# Physical or OS indexes

- obj->os\_index
  - The ID given by the OS/hardware
- P#3
  - Default in Istopo graphic mode
  - Istopo -p
- NON PORTABLE
  - Depend on motherboards, BIOS, version, ...
- DON'T USE THEM

# Logical indexes

- obj->logical\_index
  - The index among an entire level
- L#2
  - Default in Istopo except in graphic mode
  - Istopo -l
- Always represent proximity (depth-first walk)
- PORTABLE
  - Does not depend on OS/BIOS/weather
- That's what you want to use

# But I still need OS indexes when binding ?!

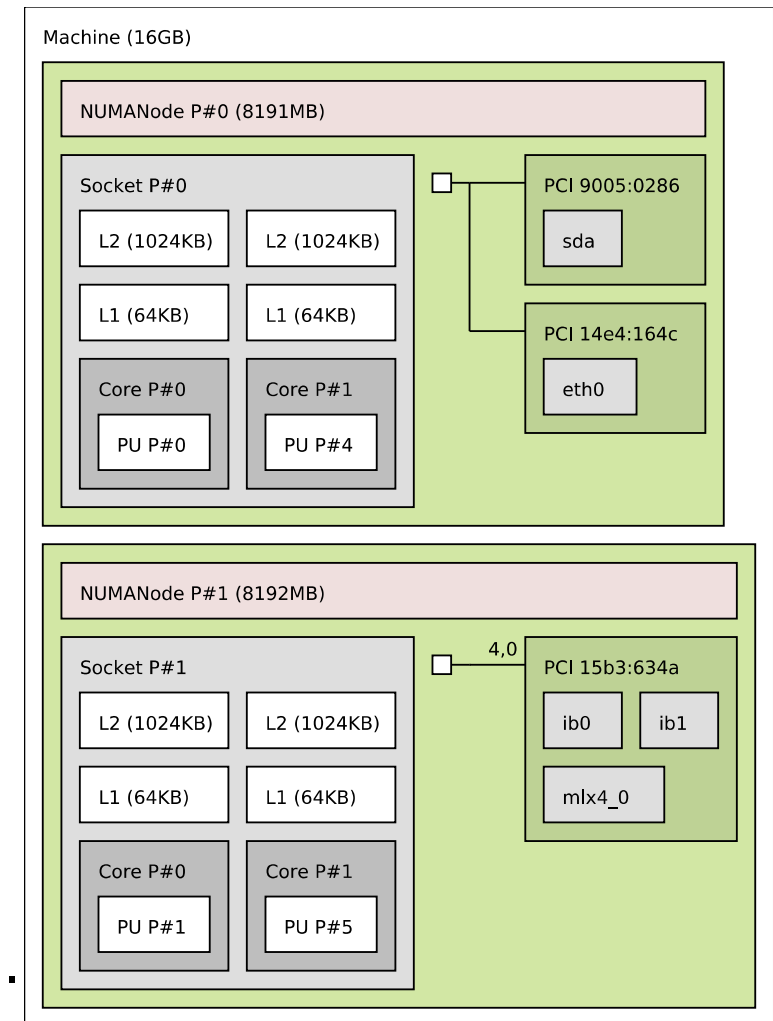
- NO !
- Just use hwloc for binding, you won't need physical/OS indexes ever again
- If you want to bind the execution to a core
  - hwloc\_set\_cpupbind(core->cpuset)
  - Other API functions for binding entire processes, single thread, memory, for allocating bound memory, etc.

# Bitmap, CPU sets, Node sets

- Generic mask of bits : `hwloc_bitmap_t`
  - Possibly infinite
  - Opaque, used to describe object contents
    - Which PU are inside this object (`obj->cpuset`)
    - Which NUMA nodes are close to this object (`obj->nodeset`)
  - Can be combined to bind to multiple cores, etc.
    - and, or, xor, not, ...

# I/O devices

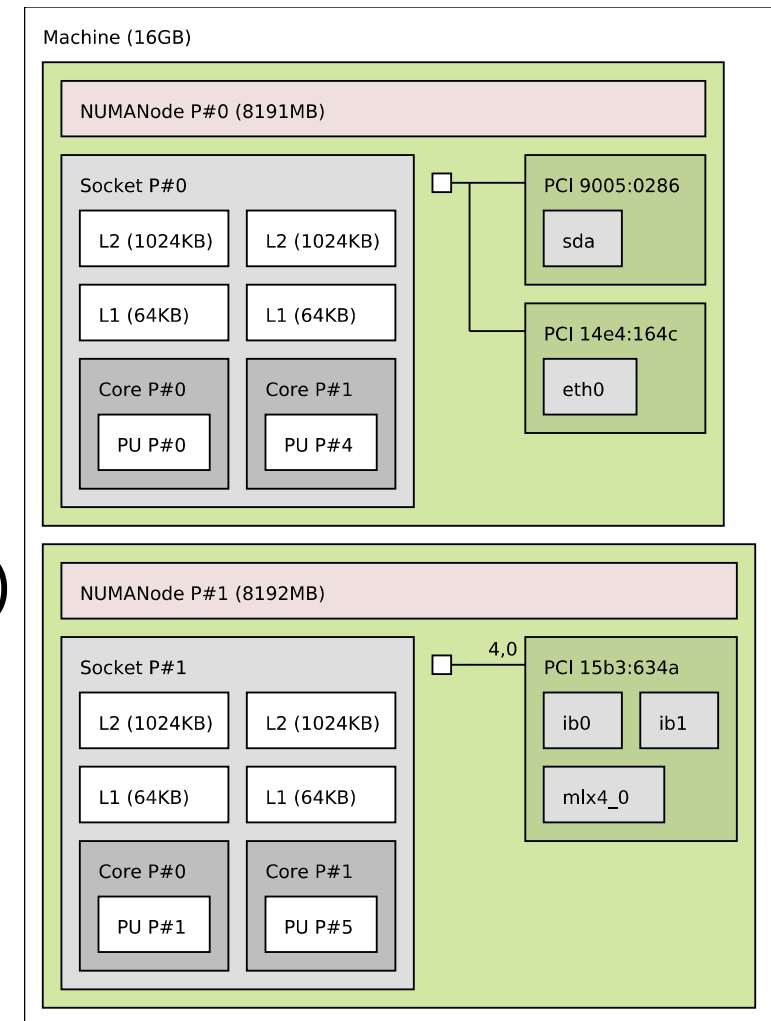
- Binding tasks near the devices they use improves their data transfer time
  - GPUs, high-performance NICs, InfiniBand, ...
- You cannot bind tasks or memory on these devices
  - But these devices may have interesting attributes
    - Device type, GPU capabilities, embedded memory, link speed, ...





# I/O objects

- Some I/O trees are attached to the object they are close to
- PCI device objects
  - Optional I/O bridge objects
- How to match your software handle with a PCI device ?
  - OS/Software devices (when known)
    - sda, eth0, ib0, mlx4\_0
- Disabled by default
  - Except in Istopo



# Extended attributes

- obj->userdata pointer
  - Your application may store whatever it needs there
  - hwloc won't look at it, it doesn't know what's it contains
- (name,value) info attributes
  - Basic string annotations, hwloc adds some
    - HostName, Kernel Release, CPU Model, PCI Vendor, ...
  - You may add more

# Configuring the topology

- Between `hwloc_topology_init()` and `load()`
  - `hwloc_topology_set_xml()`, `set_synthetic()`
  - `hwloc_topology_set_flags()`, `set_pid()`
  - `hwloc_topology_ignore_type()`
- After `hwloc_topology_load()`
  - `hwloc_topology_restrict()`
  - `hwloc_topology_insert_misc_object...`

# Helpers

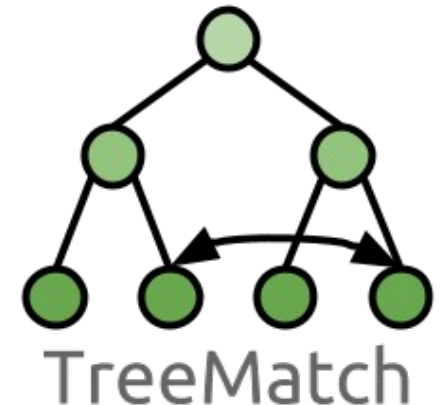
- hwloc/helpers.h contains a lot of helper functions
  - Iterators on levels, children, restricted levels
  - Finding caches
  - Converting between cpusets and nodesets
  - Finding I/O objects
  - And much more
- Use them to avoid rewriting basic functions
- Use them to understand how things work and write what you need

# 7

## Use cases

# MPI process placement

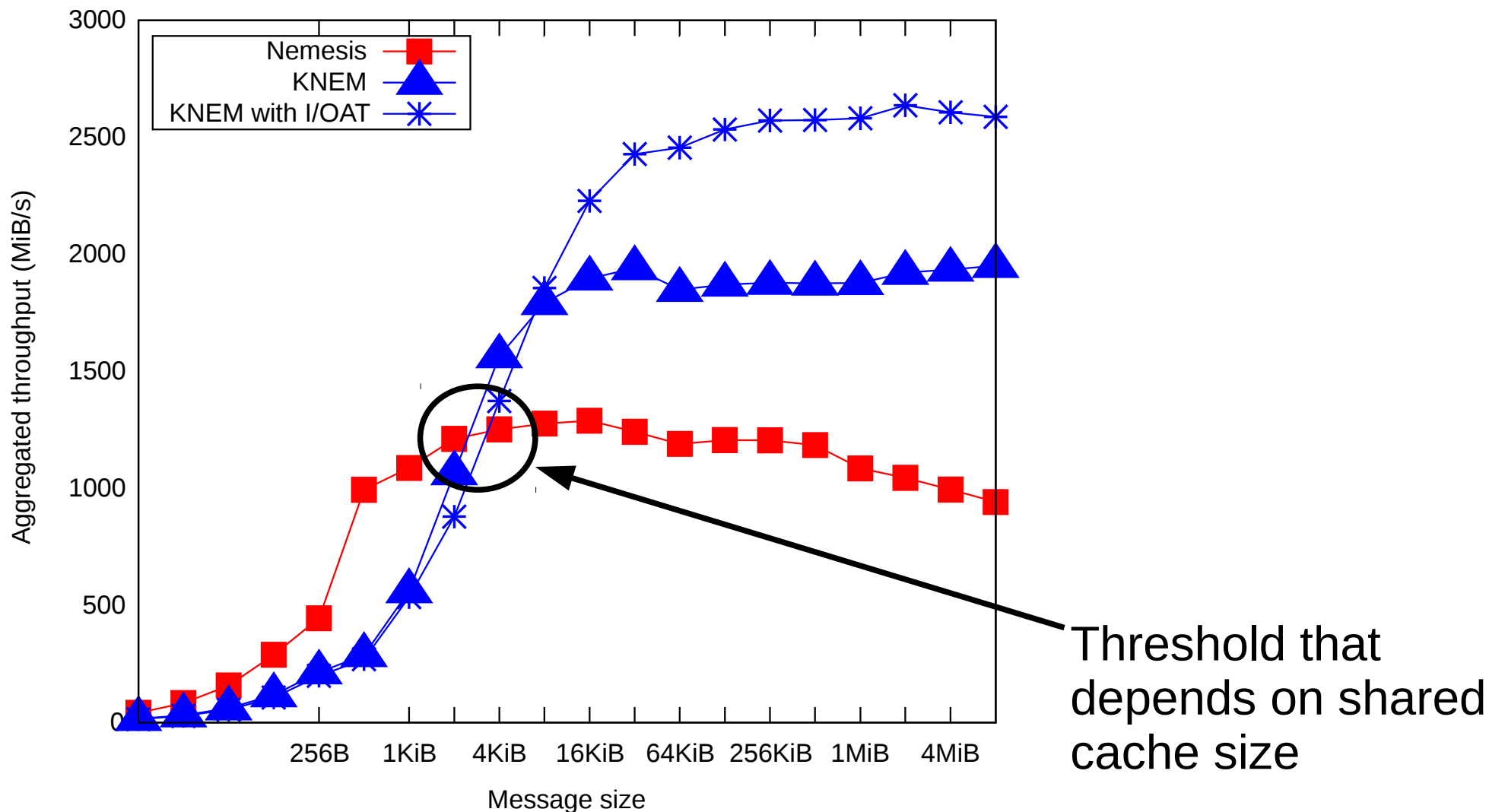
- Given a matrix describing the communication pattern of an application
- How to place processes communicating intensively on nearby cores ?
- This becomes a mapping of a tree of processes
  - Ordered by communication intensiveness
- ... onto a tree of hardware resources
  - Given by hwloc



# OpenMP thread scheduling with ForestGOMP

- OpenMP threads of the same parallel section often needs fast synchronization
  - Must stay together on the machine
    - Shared caches improve synchronization
- Build a tree of OpenMP teams and threads
  - Grouped by software affinities
- ... and map it onto a tree of hardware caches, cores, NUMA nodes, ...
  - Grouped by hardware locality

# Topology-aware thresholds for MPI intra-node communication





# Advanced binding strategies in MPI

- Open MPI
  - `mpiexec --bind-to core --map-by core ...`
    - Map by node
  - `mpiexec --bind-to core --mca rmaps_lama_map nsc`
    - Map by node, then by socket, then by core
  - See `mpiexec --help`
- MPICH
  - `mpiexec -bind-to core -map-by BSC ...`
    - Map by node (board), then by socket, then by core
  - See `mpiexec -bind-to -help`

# What about OpenMP ?

- Still far from MPI
  - Both for features and for portability of options

# 8

## Conclusion

# More information

- The documentation
  - <http://www.open-mpi.org/projects/hwloc/doc/>
  - Related pages
    - <http://www.open-mpi.org/projects/hwloc/doc/v1.8/pages.php>
  - FAQ
    - <http://www.open-mpi.org/projects/hwloc/doc/v1.8/a00028.php>
- 3-4 hours tutorials with exercises on the webpage
- README and HACKING in the source
- [hwloc-users@open-mpi.org](mailto:hwloc-users@open-mpi.org) for questions
- [hwloc-devel@open-mpi.org](mailto:hwloc-devel@open-mpi.org) for contributing
- [hwloc-announce@open-mpi.org](mailto:hwloc-announce@open-mpi.org) for new releases
- <https://svn.open-mpi.org/trac/hwloc/report> for reporting bugs

Thanks!

Questions?



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